# Database Tuning for Better Performance by Rafi Balbirsky

*In association with:* 

PerformanceArt



#### Webinar Targets

- Present the challenges of managing performance in Multi-DB environments
  (Oracle, SQL Server, MySQL, DB2)
- Focus on the common performance issues of all DB applications and establish a common language
- Introduce a new methodology for tuning applications across different databases

#### Multi-DB Challenges

- Business processes and applications are spanned across different DBs
- Each database will have its own set of performance tools and guidelines for best practice – and its own team of internal and external advisors

# Multi-DB Challenges

Business Performance information is fragmented over different databases

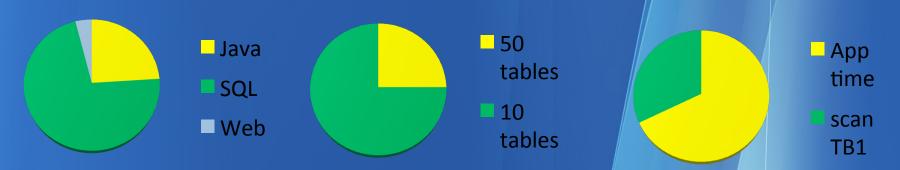
#### Performance of DB Applications

- 80% of application performance problems lie in the application
- Dealing with them by DB, storage or hardware upgrades is costly, not scalable and has very poor ROI compared to application tuning

#### Pareto Principle in SQL Tuning

- 80% of application time is spent in 20% of the application code (SQL calls)
- > 80% of application DB time is spent in 20% of the SQL statements it issues
- 80% of SQL time is spent in 20% of its Execution Plan steps

#### Application time breakdown



#### Different Databases – Same Issues

- Oracle, SQL server, MySQL and DB2 are different in many aspects but they use B-tree indices in basically the same way.
- They all suffer in the same way from sort on disk, heavy joins, full scans in accessing the data and index overhead in updating the data

# Application Time in DB Breakdown – Typical Appearance



# Resource Pie – Optimal Appearance



#### Common Performance Issues

#### Application layer

- Rerunning SQL
- Reparsing SQL

The solution is in the application.

#### Common Performance Issues

#### **DB** Layer

- Full scans
- Index range
- Index overhead
- Sorts
- Joins

#### New Approach

- A methodology for tuning objects by schema which is applicable to all DB types
- Enables a unified view of performance across all databases

# Key Metrics of Object Performance

- Access time
- Access type (full scan, index range, sort, join etc.)
- Access wait (IO, CPU, Lock, DB engine)
- Statistics (Rows#, distinct values etc.)

# Traditional Application Tuning

Currently, DB tuning and monitoring efforts use two methodologies as a starting point for the tuning process:

- Instance level Instance stats
- SQL statement level top statements

#### What's Wrong with That?

- Poor link between schema / business elements and performance information
- Poor root cause analysis
- Poor visibility to actual usage of schema objects
- Missing highly beneficial tuning opportunities

# Top SQL is Misleading

Action	ns Schedule SQL Tuning Advisor	<b>y</b> Go	
Select All   Select None			
Select	Activity (%) ▽	SQL ID	SQL Type
	89.70	05b6pvb81dg8b	SELECT
	1.52	5sd5bg2mgt4dj	SELECT
	I 1.21	2b064ybzkwf1y	PL/SQL EXECUTE
	.91	cxjqbfn0d3yqq	SELECT
	.61	g17xjfym83gqb	SELECT
	.61	05xcf43d9psvm	SELECT
	.30	6htq3p9j91y0s	INSERT
	.30	05xcf43d9psvm	SELECT
	.30	4y6zhv13yf146	SELECT
	.30	51vw8qf5uprrv	SELECT
Action	ns Schedule SQL Tuning Advisor	<b>∨</b> (Go)	

#### Real Application Bottlenecks

Many statements use the same access patterns on the same objects but won't show up in the top SQL.

Their aggregated impact is missed, together with their tuning opportunities.

CPU 🛭 WAIT 2 10 TOTAL 2 AME 📳 SQL\_ID SQL\_TEXT 577 815 SELECT COUNT(\*) FROM PAYMNT.PAYMENTS WHERE TRUNC(PAYMENT TIMESTAMP) = TRUNC(:b1) 🛝 qya73as76zut0 235 g9rh4b2h9x1nj 251 2 562 815 SELECT COUNT(\*) FROM PAYMNT.PAYMENTS\_WHERE TRUNC(PAYMENT\_TIMESTAMP) = TRUNC(:b1) 👍 9a1zvvg3kgama 251 3 559 813 SELECT COUNT(\*) FROM PAYMNT.PAYMENTS\_WHERE TRUNC(PAYMENT\_TIMESTAMP) = TRUNC(:b1) / 6vnuy4xcabq9u 242 2 569 813 SELECT COUNT(\*) FROM PAYMNT.PAYMENTS\_WHERE TRUNC(PAYMENT\_TIMESTAMP) = TRUNC(:b1) / 564 amjnjpgr2zypk 3 810 SELECT COUNT(\*) FROM PAYMNT.PAYMENTS\_WHERE TRUNC(PAYMENT\_TIMESTAMP) = TRUNC(:b1) / 243 6hy0c8y1d45y5 1 0 643 644 SELECT PAYMENT\_CODE FROM PAYMENTS WHERE PHONE = TRIM(:B2 ) AND PHONE\_PRE = TRIM(:B1 ) A 1s4666kcuzhdb 127 0 0 127 SELECT COUNT(\*) FROM PAYMNT.PAYMENTS WHERE TRUNC(PAYMENT\_TIMESTAMP) = TRUNC(:b1) / 87 98 SELECT ROVID, PAYMENT\_CODE, PAYMENT\_TIMESTAMP, PAYMENT\_BY, CUST, PAYMENT\_TYPE, FORM\_N c5fzf1whpf7rd 10 1 86cr03g23p46x 0 0 64 64 SELECT PAYMENT\_CODE,S.CUST,PAYMENT\_TYPE, GET\_CUST\_NAME(S.CUST) CUST\_NAME, REPLACE(

27 SELECT PAYMENT\_CODE FROM PAYMENTS WHERE ID\_NUMBER = TRIM(:B1 ) AND PAYMENT\_TIMESTAM

16 SELECT PAYMENT CODE FROM PAYMENTS WHERE PAYMENT CODE NOT IN (SELECT PAYMENT CODE

16 SELECT ROWID, PAYMENT\_CODE, PAYMENT\_TIMESTAMP, PAYMENT\_BY, CUST, PAYMENT\_TYPE, FORM\_N

11 SELECT COUNT(\*) FROM PAYMNT.PAYMENTS WHERE PAYMENT\_TIMESTAMP BETWEEN TO\_DATE(:b1

11 SELECT SUM(CALLS\_ANSWERED) FROM CTI.CTI\_ICDNSTAT WHERE CDN=: B2 AND TRUNC(TIMESTAMP)

5 DECLARE job BINARY\_INTEGER := :job; next\_date DATE := :mydate;\_broken BOOLEAN := FALSE; BEGIN d

4 select p.cust,count(\*), (select count(\*) from payments p1 where TRUNC(P1.PAYMENT\_TIMESTAMP) =TRU

4 insert into wri\$\_adv\_objspace\_trend\_data select timepoint, space\_usage, space\_alloc, quality from table

4 SELECT NVL(TRUNC(SUM(CALL\_LENGTH) / 60 ,2),0) FROM PAYMNT.PAYMENTS WHERE PAYMENT\_TI

2 SELECT U.SPACE\_USED, U.SPACE\_ALLOCATED FROM TABLE(DBMS\_SPACE.OBJECT\_SPACE\_USAGE\_1

2 SELECT ROWID, PAYMENT\_CODE, PAYMENT\_TIMESTAMP, PAYMENT\_BY, CUST, PAYMENT\_TYPE, FORM\_N

Line 54 Column 45

...סילבוס לטלאו 📶

2 insert into payment payments (payment code payment timestamp payment by payment at cust payment 🚩

Findings.pptx

| Modified | Windows: CR/... Ed

< 🔚 😻 📗 🕕

tput 闍 Explain 🕍 Autotrace 📵 DBMS Output 🞧 OWA Output

0

3

1

1

0

4

1

4 0

0

0

11

27

13

15

10

0

0

0

3

0

2

2

🕊 EditPlus - [C:\D...

0

0

0

0

0

5

0

0

0

0

0

0

C:\Documents ...

5zs50mjha9n74

7w78fpu7ntcp8

bwaps49vt024h 7h5v67mc6fzi0

bg8x29vhbn9cs

9nnzz1jn5f41s

86mm6ythqfbqz

2n6xk782a8ray

7wqks43wrjtrz

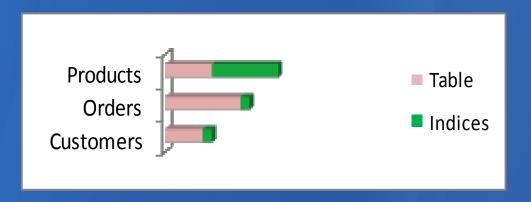
8ay8h4m2wcbdi

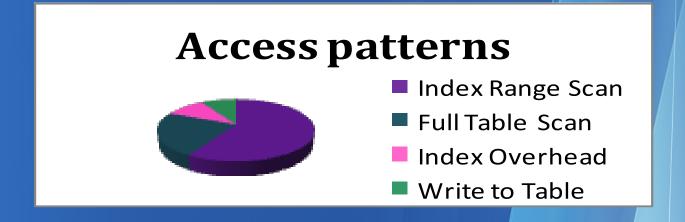
06dnm2r3vkbqv

C:\WINDOWS\...

8szmwam7fysa3

#### Aggregate by Application Objects





#### Object Tuning Methodology

#### Anchor Point – Top Objects

- Start with the central and most used tables (scope can be instance or application)
- Table *total* access time includes all the access time of its indices

# Create Tuning Sets from Explain Plan and Performance Information

A tuning set consists of SQL statements grouped by:

- Application identifiers
- Application objects
- Access types

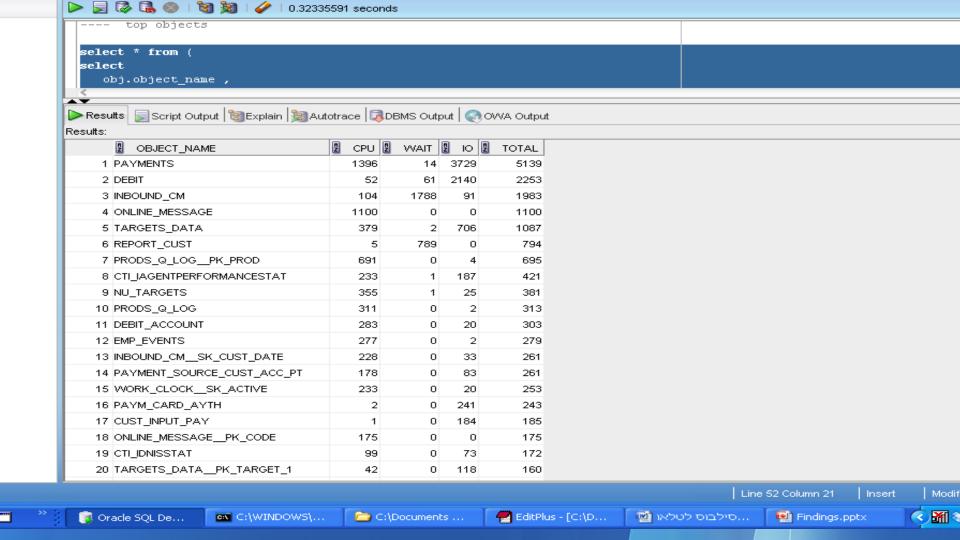
#### **Object Tuning Methodology**

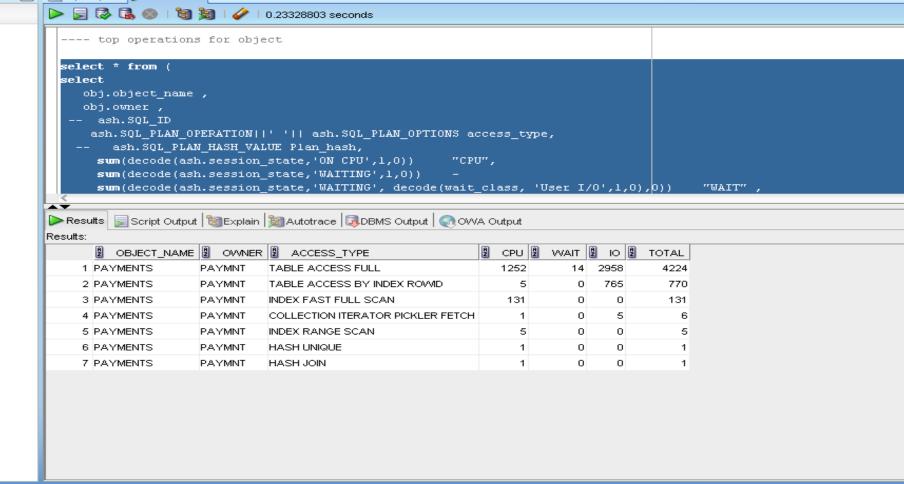
#### Applies for all relational DB:

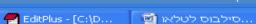
- Full table scan has different names in different DB but means the same
- Range scans on B-tree indices in all DBs have the same inefficiencies
- Adding indices to a table will always degrade update activities
- Sorts in all DBs use temporary space on disk when the sort result set is too big to be handled in memory

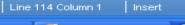
# Aggregate Full Scan Time on a Table

- See how much time the application is spending on full scans
- Look at the top SQL in this group and identify possible indices













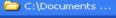




































CPU 🛭 WAIT 2 10 TOTAL 2 AME 📳 SQL\_ID SQL\_TEXT 577 815 SELECT COUNT(\*) FROM PAYMNT.PAYMENTS WHERE TRUNC(PAYMENT TIMESTAMP) = TRUNC(:b1) 🛝 qya73as76zut0 235 g9rh4b2h9x1nj 251 2 562 815 SELECT COUNT(\*) FROM PAYMNT.PAYMENTS\_WHERE TRUNC(PAYMENT\_TIMESTAMP) = TRUNC(:b1) 👍 9a1zvvg3kgama 251 3 559 813 SELECT COUNT(\*) FROM PAYMNT.PAYMENTS\_WHERE TRUNC(PAYMENT\_TIMESTAMP) = TRUNC(:b1) / 6vnuy4xcabq9u 242 2 569 813 SELECT COUNT(\*) FROM PAYMNT.PAYMENTS\_WHERE TRUNC(PAYMENT\_TIMESTAMP) = TRUNC(:b1) / 564 amjnjpgr2zypk 3 810 SELECT COUNT(\*) FROM PAYMNT.PAYMENTS\_WHERE TRUNC(PAYMENT\_TIMESTAMP) = TRUNC(:b1) / 243 6hy0c8y1d45y5 1 0 643 644 SELECT PAYMENT\_CODE FROM PAYMENTS WHERE PHONE = TRIM(:B2 ) AND PHONE\_PRE = TRIM(:B1 ) A 1s4666kcuzhdb 127 0 0 127 SELECT COUNT(\*) FROM PAYMNT.PAYMENTS WHERE TRUNC(PAYMENT\_TIMESTAMP) = TRUNC(:b1) / 87 98 SELECT ROVID, PAYMENT\_CODE, PAYMENT\_TIMESTAMP, PAYMENT\_BY, CUST, PAYMENT\_TYPE, FORM\_N c5fzf1whpf7rd 10 1 86cr03g23p46x 0 0 64 64 SELECT PAYMENT\_CODE,S.CUST,PAYMENT\_TYPE, GET\_CUST\_NAME(S.CUST) CUST\_NAME, REPLACE(

27 SELECT PAYMENT\_CODE FROM PAYMENTS WHERE ID\_NUMBER = TRIM(:B1 ) AND PAYMENT\_TIMESTAM

16 SELECT PAYMENT CODE FROM PAYMENTS WHERE PAYMENT CODE NOT IN (SELECT PAYMENT CODE

16 SELECT ROWID, PAYMENT\_CODE, PAYMENT\_TIMESTAMP, PAYMENT\_BY, CUST, PAYMENT\_TYPE, FORM\_N

11 SELECT COUNT(\*) FROM PAYMNT.PAYMENTS WHERE PAYMENT\_TIMESTAMP BETWEEN TO\_DATE(:b1

11 SELECT SUM(CALLS\_ANSWERED) FROM CTI.CTI\_ICDNSTAT WHERE CDN=: B2 AND TRUNC(TIMESTAMP)

5 DECLARE job BINARY\_INTEGER := :job; next\_date DATE := :mydate;\_broken BOOLEAN := FALSE; BEGIN d

4 select p.cust,count(\*), (select count(\*) from payments p1 where TRUNC(P1.PAYMENT\_TIMESTAMP) =TRU

4 insert into wri\$\_adv\_objspace\_trend\_data select timepoint, space\_usage, space\_alloc, quality from table

4 SELECT NVL(TRUNC(SUM(CALL\_LENGTH) / 60 ,2),0) FROM PAYMNT.PAYMENTS WHERE PAYMENT\_TI

2 SELECT U.SPACE\_USED, U.SPACE\_ALLOCATED FROM TABLE(DBMS\_SPACE.OBJECT\_SPACE\_USAGE\_1

2 SELECT ROWID, PAYMENT\_CODE, PAYMENT\_TIMESTAMP, PAYMENT\_BY, CUST, PAYMENT\_TYPE, FORM\_N

Line 54 Column 45

...סילבוס לטלאו 📶

2 insert into payment payments (payment code payment timestamp payment by payment at cust payment 🚩

Findings.pptx

| Modified | Windows: CR/... Ed

< 🔚 😻 📗 🕕

tput 闍 Explain 🕍 Autotrace 📵 DBMS Output 🞧 OWA Output

0

3

1

1

0

4

1

4 0

0

0

11

27

13

15

10

0

0

0

3

0

2

2

🕊 EditPlus - [C:\D...

0

0

0

0

0

5

0

0

0

0

0

0

C:\Documents ...

5zs50mjha9n74

7w78fpu7ntcp8

bwaps49vt024h 7h5v67mc6fzi0

bg8x29vhbn9cs

9nnzz1jn5f41s

86mm6ythqfbqz

2n6xk782a8ray

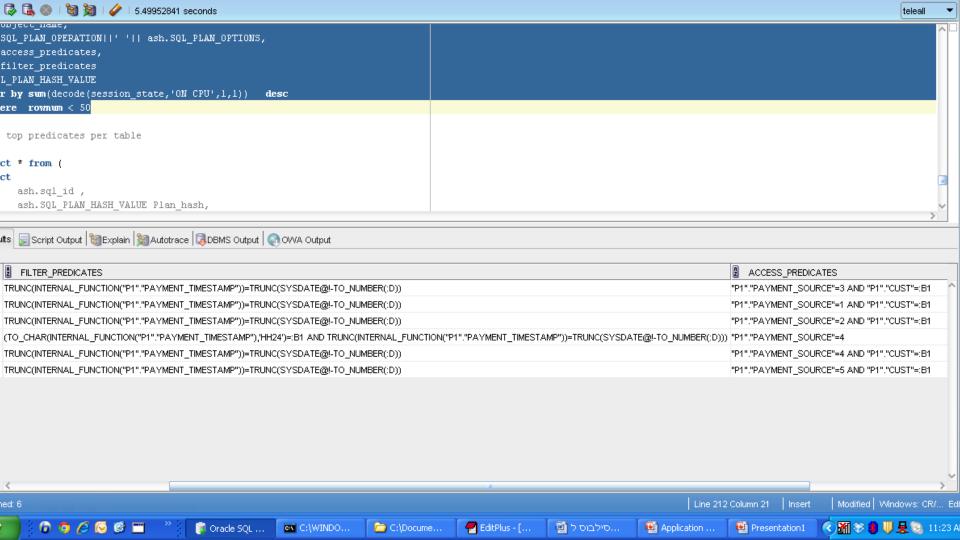
7wqks43writrz

8ay8h4m2wcbdi

06dnm2r3vkbqv

C:\WINDOWS\...

8szmwam7fysa3



#### **B-tree Indices**

- Most applications rely heavily on B-tree indices across all DBs
- Indices on large and central tables can be large and consume high amounts of I/O, CPU and memory
- Avoiding full scans by using indices can create new performance issues, slowing down application responsiveness

# Aggregate Index Time to Table

#### Now you can assess:

- Total time of access to a table
- Index usage
- Index misusage
- Index overhead

Now you can optimize your index structure

#### Object Issues – Bad Index Structure

#### Common index related problems:

- Matching level / Clustering Factor
- Index overhead
- Change to index structure can boost performance
- Tuning gain can be calculated

# Index Only / Covering Index Access

- No table blocks are read
- Clustering factor has no effect
- Sort can be avoided

#### Index Overhead

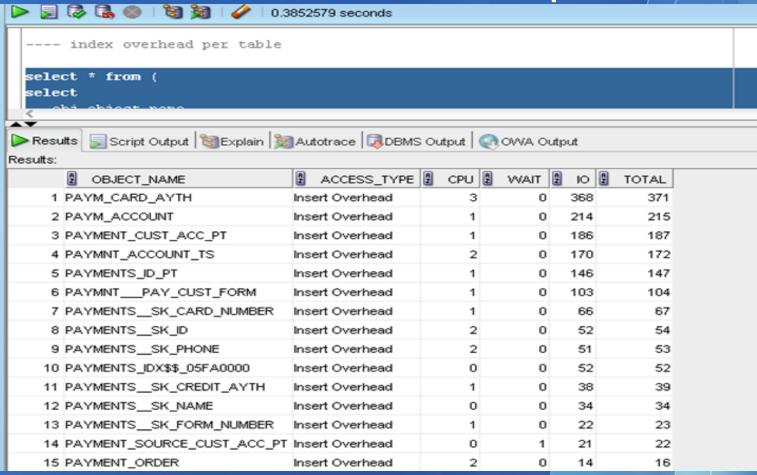
Finally get an answer for this important question

How much do I pay for my indices?

# Index Overhead Solution

- Detect unused indices
- Detect subset indices
- Detect unused columns that inflate index size

#### Index Overhead – Real Example



#### Schema Analysis - Local

- Often more effective than SQL tuning
- Solid entity for DBA. Fits well in solution space
- Helps in optimizing the match between application design and usage
- Exposes hidden performance bottlenecks
- Reveals usage of schema objects
- Enables effective usage of DB advisors

#### Schema Analysis - Global

- Enables continuous and full visibility of DB layer impact on business performance
- Cuts down the time and cost of performance management
- Increases company's scalability and its flexibility in answering business needs

#### Thank you for your attention

Questions?

